Ex2:

$$S1 = \{A/?x, B/?y, C/?w, D/?z\}$$

 $S2 = \{g(?x, ?y)/?z\}$
 $S1S2 = \{A/?x, B/?y, C/?w, D/?z\}$

Change the order of the substitutions in the previous example.

 $S2S1 = \{ A/?x, B/?y, C/?w, g(A, B)/?z \}$ $S1S2 \neq S2S1$

Verifying Consistency of Substitutions

To determine consistency, create two sets, T, V. Put all terms for all involved substitutions in T, and all variables for all involved substitutions in V. Then propagate values of the variables. If any variable has to take more than one constant value, then substitutions are inconsistent. Ex: We have the following substitutions:



How do we find substitutions?

$$\frac{W1(?x) \Rightarrow W2(?x)}{W1(A)} \} \Rightarrow W2(A)$$

Use the recursive Unify procedure for unifying formulas (wffs) and finding substitutions.

Recursive Procedure **UNIFY**(E1, E2)

- 1 **if** either E1 or E2 is an atom (that is, a predicate symbol, a function symbol, a constant symbol, a negation symbol or a variable), interchange the arguments E1 and E2 (is necessary) so that E1 is an atom, and **do:**
- begin 2 3 if E1 and E2 are identical, return NIL if E1 is a variable, do: 4 begin 5 if E1 occurs in E2, return FAIL 6 7 **return** {E2/E1} 8 end 9 if E2 is a variable, return {E1/E2} 10 return FAIL 11 end contd...

...contd

- 12 F1 <- the first element of E1, T1 <- the rest of E1
- 13 F2 <- the first element of E2, T2 <- the rest of E2
- 14 Z1 <- **UNIFY**(F1,F2)
- 15 **if** Z1 = FAIL, **return** FAIL
- 16 G1 <- result of applying Z1 to T1
- 17 G2 <- result of applying Z1 to T2
- 18 Z2 <- **UNIFY**(G1, G2)
- 19 **if** Z2 = FAIL, **return** FAIL
- 20 return the composition of Z1 and Z2

WE

<u>Ex:</u> {P[A, ?x, f(?y)], P[?y, B, f(?z)]}

First convert everything into lists.

E1 = (P A ?x (f ?y)) E2 = (P ?y B (f ?z))

- F1 = P T1 = (A ?x (f ?y))
- F2 = P T2 = (?y B (f ?z))

 $Z1 \leftarrow UNIFY(F1, F2) \rightarrow return NIL$

$$G1 = T1 = (A ?x (f ?y))$$

 $G2 = T2 = (?y B (f ?z))$

$Z2 \leftarrow UNIFY ((A ?x (f ?y)) (?y B (f ?z)))$ $E1 \qquad E2$

F1 = A	T1 = (?x (f ?y))
F2 = ?y	T2 = (B (f ?z))

Z1 \leftarrow UNIFY (A, ?y) \rightarrow return (A/?y) F1 F2

G1 = (?x (f A))G2 = (B (f ?z))

$Z2 \leftarrow UNIFY ((?x (f A)) (B (f ?z)))$ E1 E2

F1 = ?xT1 = ((f A))F2 = BT2 = ((f ?z))

 $Z1 \leftarrow UNIFY(?x B) \rightarrow return (B/?x)$

 $Z2 \leftarrow UNIFY(((f A)) ((f ?z)))$: $Z2 \leftarrow UNIFY((A) (?z))$

F1 = AT1 = ()F2 = ?zT2 = ()

Z1 \leftarrow UNIFY (A ?z) \rightarrow return (A/?z)

Compose
$$((A/?z) (B/?x))$$

 \downarrow
 $((A/?z, B/?x))$

Compose
$$((A/?y) (A/?z, B/?x))$$

 $((A/?y, A/?z, B/?x))$

Compose (NIL (A/?y, A/?z, B/?x))

Rule Systems: Simplest Knowledge Systems

$$\begin{aligned} & \text{Rules} \begin{cases} W1(v_1, v_2...) => W1'(...) \\ W2(...) => W2'(...) \\ \vdots \\ Wn(...) => Wn'(...) \\ S_i = \{C_1/v_1, C_2/v_2, ..\} \\ \text{generated} \\ \text{fact} \end{cases} \\ & \text{Facts} \begin{cases} W_{f1}(C1, C2, ...) \\ W_{f2}(C1, C2, ...) \\ \vdots \end{cases} \end{aligned}$$

Forward Chaining: match fact expressions with antecedents, generate new facts.

Backward Chaining: match goal or subgoal expressions with consequents, generate new subgoals.

procedure Forward-Chain (*KB*, *p*)

if there is a sentence in *KB* that is a renaming of *p* then return

```
Add p to KB

for each in KB such that for some I,

UNIFY (p_i, p) = \alpha t succeeds do

FIND-AND-INFER (KB, [p_i, ..., p_{i-1}, p_{i+1}, ..., p_n], q, \theta)

end
```

```
procedure FIND-AND-INFER (KB, premises,
```

conclusions, θ)

```
if premises = [] then

FORWARD-CHAIN (KB, SUBST (\theta, conclusion))

else for each p' in KB such that UNIFY (p', SUBST(\theta,

FIRST (premises))) = \theta_2 do

FIND-AND-INFER (KB, REST (premises), conclusion,

COMPOSER (\theta, \theta_2))
```

end

Figure - The forward-chaining inference algorithm. It adds to *KB* all the sentences that can be inferred from the sentence *p*. If *p* is already in *KB*, it does nothing. If *p* is new, consider each implication that has a premise that matches *p*. For each such implication, if all the remaining premises are in *KB*, then infer the conclusion. If the premises can be matched several ways, then infer each corresponding conclusion. The substitution of θ keeps track of the way things match.

Review and Exercises

• The following expressions are given:

Manager(PURCHASING-DEPT, JOHN-JONES) Works-in(PURCHASING-DEPT, JOE-SMITH)

We also have the following production:

[Works $-in(?x, ?y) \wedge Manager(?x, ?z)$] $\Rightarrow Boss - of(?y, ?z)$

- a) Using these expressions, show the forward inference steps which show that JOHN-JONES is the boss of JOE-SMITH.
- b) Write a goal for finding the boss of JOE-SMITH. Using the above expressions, by setting up subgoals through backward inference steps and matching show JOHN-JONES is his boss.







1b) goal: Boss-of(JOE-SMITH, ?t) S2 = {JOE-SMITH /?y, ?t /?z}

New subgoals: {Works-in(?x, ?y) S2 Manager(?x, ?z) S2 {Works-in(?x, JOE-SMITH) sg1 Manager(?x, ?t) sg2

Match with facts in database: $S3 = \{PURCHASING-DEPT/?x\}$ Works-in(PURCHASING-DEPT, JOE-SMITH) $sg1 S3 = W_{f2}$

 $S4 = \{PURCHASING-DEPT/?x, JOHN-JONES /?t\}$ $sg2 \ S4 = W_{f1}$



